

14

$$C_k \equiv M_k^{e_k} \pmod{p_k},$$

15 where

$$M_1 \equiv M \pmod{p_1},$$

$$M_2 \equiv M \pmod{p_2},$$

⋮

$$M_k \equiv M \pmod{p_k},$$

20

$$e_1 \equiv e \pmod{p_1 - 1},$$

$$e_2 \equiv e \pmod{p_2 - 1}, \text{ and}$$

⋮

$$e_k \equiv e \pmod{p_k - 1},$$

25

26 where  $e$  is a number relatively prime to  $(p_1 - 1), (p_2 - 1), \dots, (p_k - 1)$ ,

27 solving said subtasks to determine results  $C_1, C_2 \dots C_k$ ,

28 combining said results of said subtasks in accordance with a fast recursive combining

29 process to produce said ciphertext word signal  $C$  whereby,

$$30 Y_i \equiv Y_{i-1} + [(C_i - Y_{i-1}) (w_i^{-1} \pmod{p_i}) \pmod{p_i}] \cdot w_i \pmod{n}$$

31  $2 \leq i \leq k$ , and

$$32 C = Y_k, Y_1 = C_1, \text{ and } w_i = \prod_{j < i} p_j$$

33 whereby processing of a minimal amount of computer instructions is required for said  
34 step of encoding.

1 15. (Twice Amended) A method for establishing cryptographic communications that are  
2 backwards compatible with preexisting public key infrastructures, comprising the steps of:  
3 decoding a ciphertext word  $C$  to a message word  $M$ , wherein  $M$  corresponds to a number  
4 representative of a message and wherein,

$$5 0 \leq M \leq n-1$$

6 wherein  $n$  is a composite number formed by the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ ,  $k$  is an integer greater  
7 than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers,  $C$  is a number representative of an

8 encoded form of message word M that is encoded by transforming said message word M to said  
9 ciphertext word C whereby,

10  $C \equiv M^e \pmod{n}$ ,

11 and wherein e is a number relatively prime to  $(p_1-1)$ ,  $(p_2-1)$ , ..., and  $(p_k-1)$ ,

12 said decoding step being performed using a decryption exponent d that is defined by

13  $d \equiv e^{-1} \pmod{(p_1-1)(p_2-1) \dots (p_k-1)}$ ,

14 said decoding step including the steps of,

15 (i) defining a plurality of k sub-tasks in accordance with

16  $M_1 \equiv C_1^{d_1} \pmod{p_1}$ ,

17  $M_2 \equiv C_2^{d_2} \pmod{p_2}$ ,

18  $\vdots$

19  $M_k \equiv C_k^{d_k} \pmod{p_k}$ ,

20

21 where

22  $C_1 \equiv C \pmod{p_1}$ ,

23  $C_2 \equiv C \pmod{p_2}$ ,

24  $\vdots$

25  $C_k \equiv C \pmod{p_k}$ ,

26

27  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,

28  $d_2 \equiv d \pmod{(p_2 - 1)}$ , and

29  $\vdots$

30  $d_k \equiv d \pmod{(p_k - 1)}$ ,

31 (ii) solving said sub-tasks to determine results  $M_1, M_2, \dots, M_k$ , and

32 (iii) combining said results of said subtasks in accordance with a fast recursive combining  
33 process to produce said message word M in accordance with,

34  $Y_i \equiv Y_{i-1} + [(M_i - Y_{i-1}) (w_i^{-1} \pmod{p_i}) \pmod{p_i}] \cdot w_i \pmod{n}$

35 where  $2 \leq i \leq k$ , and

36  $M = Y_k, Y_1 = M_1, \text{ and } w_i = \prod_{j < i} p_j$

37 whereby processing of a minimal amount of computer instructions is required for said  
38 step of decoding.

1 16. (Twice Amended) A cryptographic communications system for establishing communications  
2 that are backwards compatible with preexisting public key infrastructures, comprising:

3 a communication medium;

4 [an] encoding means coupled to said communication medium and adapted for  
5 transforming a transmit message word M to a ciphertext word C and for transmitting said  
6 ciphertext word C on said medium, where M corresponds to a number representative of a  
7 message, and

8  $0 \leq M \leq n-1$  where n is a composite number of the form

9  $n = p_1 \cdot p_2 \cdot \dots \cdot p_k$ ,

10 where k is an integer greater than 2 and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers,  
11 and where C corresponds to a number representative of an enciphered form of said message, and  
12 corresponds to

13  $C \equiv M^e \pmod{n}$ ,

14 where e is a number relatively prime to  $(p_1-1), (p_2-1), \dots, (p_k-1)$ ; and

15 [a] decoding means coupled to said communication medium and adapted for receiving C  
16 via said medium and for transforming C to a receive message word M' where M' corresponds to  
17 a number representative of a deciphered form of C, said decoding means being operative to  
18 perform a decryption process using a decryption exponent d that is defined by

19  $d \equiv e^{-1} \pmod{(p_1-1)(p_2-1) \dots (p_k-1)}$ ,

20 said decryption process including the steps of

21 (i) defining a plurality of k sub-tasks in accordance with,

22  $C_1 \equiv C \pmod{p_1}$ ,

23  $C_2 \equiv C \pmod{p_2}$ ,

24  $\vdots$

25  $C_k \equiv C \pmod{p_k}$ ,

26 where,

27  $d_1 \equiv d \pmod{(p_1 - 1)},$   
28  $d_2 \equiv d \pmod{(p_2 - 1)},$   
29  $\vdots$   
30  $d_k \equiv d \pmod{(p_k - 1)},$   
31  
32  $M_1' \equiv C_1^{d_1} \pmod{p_1},$   
33  $M_2' \equiv C_2^{d_2} \pmod{p_2}, \text{ and}$   
34  $\vdots$   
35  $M_k' \equiv C_k^{d_k} \pmod{p_k},$

36 (ii) solving said sub-tasks to determine results  $M_1', M_2', \dots, M_k'$ , and

37 (iii) combining said results of said subtasks by a fast recursive combining process to  
38 produce said receive message word  $M'$  in accordance with

39 
$$Y_i \equiv Y_{i-1} + [(M_i' - Y_{i-1}) (w_i^{-1} \pmod{p_i}) \pmod{p_i}] \cdot w_i \pmod{n}$$

40 where  $2 \leq i \leq k$  and

41 
$$M' = Y_k, Y_1 = M_1, \text{ and } w_i = \prod_{j < i} p_j,$$

42 [whereby] wherein  $M' = M$ .

1 17. (Once Amended) A method for establishing cryptographic communications that are  
2 backwards compatible with preexisting public key infrastructures, comprising the steps of:  
3 encoding a plaintext message word  $M$  to a ciphertext word  $C$ , wherein  $M$  corresponds to  
4 a number representative of a message and wherein

5 
$$0 \leq M \leq n-1,$$

6 wherein  $n$  is a composite number formed by the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ ,  $k$  is an integer  
7 greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers,  $C$  is a number  
8 representative of an encoded form of message word  $M$ , and wherein said encoding step  
9 comprises transforming said message word  $M$  to said ciphertext word  $C$ , whereby

10 
$$C \equiv M^e \pmod{n},$$

11 and wherein  $e$  is a number relatively prime to  $(p_1-1), (p_2-1), \dots, (p_k-1)$ ; and

12 decoding said ciphertext word C to a receive message word M', said decoding step being  
13 performed using a decryption exponent d that is defined by

14  $d \equiv e^{-1} \pmod{((p_1-1)(p_2-1)\dots(p_k-1))}$ ,

15 said decoding step including the further steps of,

16 defining a plurality of k sub-tasks in accordance with

17  $M_1' \equiv C_1^{d_1} \pmod{p_1}$ ,

18  $M_2' \equiv C_2^{d_2} \pmod{p_2}$ ,

19 :

20  $M_k' \equiv C_k^{d_k} \pmod{p_k}$ ,

21 wherein

22  $C_1 \equiv C \pmod{p_1}$ ,

23  $C_2 \equiv C \pmod{p_2}$ ,

24 :

25  $C_k \equiv C \pmod{p_k}$ ,

26 :

27  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,

28  $d_2 \equiv d \pmod{(p_2 - 1)}$ , and

29 :

30  $d_k \equiv d \pmod{(p_k - 1)}$ ,

31 solving said sub-tasks to determine results  $M_1'$ ,  $M_2'$ , ...  $M_k'$ , and

32 combining said results of said sub-tasks to produce said receive message word

33  $M'$ , [whereby] wherein  $M' = M$ .

1 22. (Once Amended) A cryptographic communications system for establishing communications  
2 that are backwards compatible with preexisting public key infrastructures, comprising:

3 a communication medium;

4 [an] encoding means coupled to said communication medium and adapted for  
5 transforming a transmit message word M to a ciphertext word C and for transmitting said

6 ciphertext word C on said medium, wherein M corresponds to a number representative of a  
7 message, and

8  $0 \leq M \leq n-1$ , wherein n is a composite number of the form,

9  $n = p_1 \cdot p_2 \cdot \dots \cdot p_k$

10 wherein k is an integer greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime  
11 numbers, and wherein said ciphertext word C corresponds to a number representative of an  
12 enciphered form of said message and corresponds to

13  $C \equiv M^e \pmod{n}$ ,

14 wherein e is a number relatively prime to  $(p_1-1), (p_2-1), \dots, (p_k-1)$ ; and

15 [a] decoding means communicatively coupled with said communication medium for  
16 receiving said ciphertext word C via said medium, said decoding means being operative to  
17 perform a decryption process for transforming said ciphertext word C to a receive message word  
18  $M'$ , wherein  $M'$  corresponds to a number representative of a deciphered form of C, said  
19 decryption process using a decryption exponent d that is defined by

20  $d \equiv e^{-1} \pmod{(p_1-1)(p_2-1) \dots (p_k-1)}$ ,

21 said decryption process including the steps of

22 defining a plurality of k sub-tasks in accordance with

23  $M_1' \equiv C_1^{d_1} \pmod{p_1}$ ,

24  $M_2' \equiv C_2^{d_2} \pmod{p_2}$ ,

25  $\vdots$

26  $M_k' \equiv C_k^{d_k} \pmod{p_k}$ ,

27 wherein

28  $C_1 \equiv C \pmod{p_1}$ ,

29  $C_2 \equiv C \pmod{p_2}$ ,

30  $\vdots$

31  $C_k \equiv C \pmod{p_k}$ ,

33  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,

34  $d_2 \equiv d \pmod{(p_2 - 1)}$ ,

35  $d_k \equiv d \pmod{(p_k - 1)},$   
36 solving said sub-tasks to determine results  $M_1', M_2', \dots, M_k'$ , and  
37 combining said results of said sub-tasks to produce said receive message word  $M'$   
38 whereby  $M' = M$ .

1 27. (Once Amended) A method for establishing cryptographic communications that are  
2 backwards compatible with preexisting public key infrastructures, comprising the step of:  
3 encoding a plaintext message word  $M$  to a ciphertext word  $C$ , wherein  $M$  corresponds to  
4 a number representative of a message, and  
5  $0 \leq M \leq n-1$ ,  
6  $n$  being a composite number formed from the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ , wherein  $k$  is an integer  
7 greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers, and wherein the ciphertext  
8 word  $C$  is a number representative of an encoded form of message word  $M$ , wherein said step of  
9 encoding includes the steps of

10 defining a plurality of  $k$  sub-tasks in accordance with

11  $C_1 \equiv M_1^{e_1} \pmod{p_1},$   
12  $C_2 \equiv M_2^{e_2} \pmod{p_2},$   
13  $\vdots$   
14  $C_k \equiv M_k^{e_k} \pmod{p_k},$

15 where

16  $M_1 \equiv M \pmod{p_1},$   
17  $M_2 \equiv M \pmod{p_2},$   
18  $\vdots$   
19  $M_k \equiv M \pmod{p_k},$   
20  
21  $e_1 \equiv e \pmod{(p_1 - 1)},$   
22  $e_2 \equiv e \pmod{(p_2 - 1)}, \text{ and}$

1 32. (Once Amended) A cryptographic communications system for establishing  
2 communications that are backwards compatible with preexisting public key infrastructures,  
3 comprising:

a communication medium;

[an] encoding means coupled to said communication medium and operative to transform a transmit message word  $M$  to a ciphertext word  $C$ , and to transmit said ciphertext word  $C$  on said medium, wherein  $M$  corresponds to a number representative of a message, and

$$0 \leq M \leq n-1.$$

9 n being a composite number formed from the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$  wherein k is an integer  
10 greater than 2[,] and  $p_1, p_2, \dots, p_k$ , are distinct random prime numbers, and wherein the ciphertext  
11 word C is a number representative of an encoded form of message word M, said encoding means  
12 being operative to transform said transmit message word M to said ciphertext word C by  
13 performing an encoding process comprising the steps of

defining a plurality of  $k$  sub-tasks in accordance with

$$C_1 \equiv M_1^{e_1} \pmod{p_1},$$

$$C_2 \equiv M_2^{e_2} \pmod{p_2},$$

•

$$C_k \equiv M_k^{e_k} \pmod{p_k},$$

19 where

$$M_1 \equiv M \pmod{p_1},$$

$$M_2 \equiv M \pmod{p_2},$$

20

23  $M_k \equiv M \pmod{p_k}$ ,

24

25  $e_1 \equiv e \pmod{(p_1 - 1)}$ ,

26  $e_2 \equiv e \pmod{(p_2 - 1)}$ , and

27  $\vdots$

28  $e_k \equiv e \pmod{(p_k - 1)}$ ,

29 wherein  $e$  is a number relatively prime to  $(p_1-1)$ ,  $(p_2-1)$ , ..., and  $(p_k-1)$ ,  
30 solving said sub-tasks to determine results  $C_1$ ,  $C_2$ , ...  $C_k$ , and  
31 combining said results of said sub-tasks to produce said ciphertext word  $C$ . (

1 37. (Once Amended) A method for establishing cryptographic communications that are  
2 backwards compatible with preexisting public key infrastructures, comprising the steps of:  
3       decoding a ciphertext word C to a message word M, wherein M corresponds to a number  
4       representative of a message and wherein  
5        $0 \leq M \leq n-1$   
6       wherein n is a composite number formed by the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ , k is an integer greater  
7       than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers, C is a number representative of an  
8       encoded form of message word M that is encoded by transforming said message word M to said  
9       ciphertext word C whereby

10         $C \equiv M^e \pmod{n}$ ,  
11        and wherein  $e$  is a number relatively prime to  $(p_1-1)$ ,  $(p_2-1)$ , ..., and  $(p_k-1)$ ;  
12        said decoding step being performed using a decryption exponent  $d$  that is defined by  
13                 $d \equiv e^{-1} \pmod{(p_1-1)(p_2-1) \dots (p_k-1)}$ ,  
14        wherein said step of decoding includes the steps of  
15                defining a plurality of  $k$  sub-tasks in accordance with

- 16  $M_1 \equiv C_1^{d_1} \pmod{p_1},$
- 17  $M_2 \equiv C_2^{d_2} \pmod{p_2},$
- 18  $\vdots$
- 19  $M_k \equiv C_k^{d_k} \pmod{p_k},$

20 wherein  
21  $C_1 \equiv C \pmod{p_1}$ ,  
22  $C_2 \equiv C \pmod{p_2}$ ,  
23  $\vdots$   
24  $C_k \equiv C \pmod{p_k}$ ,  
25  
26  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,  
27  $d_2 \equiv d \pmod{(p_2 - 1)}$ , and  
28  $\vdots$   
29  $d_k \equiv d \pmod{(p_k - 1)}$ ,  
30 solving said sub-tasks to determine results  $M_1, M_2, \dots, M_k$ , and  
31 combining said results of said sub-tasks to produce said message word  $M$ .

1 42. (Once Amended) A cryptographic communications system for establishing communications  
2 that are backwards compatible with preexisting public key infrastructures, comprising:  
3 a communication medium;  
4 [a decoding means] communicatively coupled with said communication medium for  
5 receiving a ciphertext word  $C$  via said medium, and being operative to transform said ciphertext  
6 word  $C$  to a receive message word  $M'$ , wherein a message  $M$  corresponds to a number  
7 representative of a message and wherein,  
8  $0 \leq M \leq n-1$   
9 wherein  $n$  is a composite number formed by the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ ,  $k$  is an integer greater  
10 than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers, and wherein said ciphertext word  
11  $C$  is a number representative of an encoded form of said message word  $M$  that is encoded by  
12 transforming  $M$  to said ciphertext word  $C$  whereby,

13  $C \equiv M^e \pmod{n}$ ,  
14 and wherein  $e$  is a number relatively prime to  $(p_1-1), (p_2-1), \dots, (p_k-1)$ ;  
15 said decoding means being operative to perform a decryption process using a decryption  
16 exponent  $d$  that is defined by

17  $d \equiv e^{-1} \pmod{(p_1-1)(p_2-1) \dots (p_k-1)}$ ,

18 said decryption process including the steps of  
19 defining a plurality of k sub-tasks in accordance with,

20  $M_1' \equiv C_1^{d_1} \pmod{p_1}$ ,

21  $M_2' \equiv C_2^{d_2} \pmod{p_2}$ ,

22 :

23  $M_k' \equiv C_k^{d_k} \pmod{p_k}$ ,

24 wherein,

25  $C_1 \equiv C \pmod{p_1}$ ,

26  $C_2 \equiv C \pmod{p_2}$ ,

27 :

28  $C_k \equiv C \pmod{p_k}$ ,

30  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,

31  $d_2 \equiv d \pmod{(p_2 - 1)}$ , and

32 :

33  $d_k \equiv d \pmod{(p_k - 1)}$ ,

34 solving said sub-tasks to determine results  $M_1'$ ,  $M_2'$ , ...  $M_k'$ , and

35 combining said results of said sub-tasks to produce said receive message word

36  $M'$ , whereby  $M' = M$

1 47. (Once Amended) A method for generating a digital signature comprising the step of:  
2 signing a plaintext message word  $M$  to create a signed ciphertext word  $C$ , wherein  $M$   
3 corresponds to a number representative of a message, and  
4  $0 \leq M \leq n-1$ ,

5  $n$  being a composite number formed from the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ , wherein  $k$  is an integer  
6 greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers, and wherein the signed  
7 ciphertext word  $C$  is a number representative of a signed form of message word  $M$ , wherein  
8  $C \equiv M^d \pmod{n}$ , and

9 wherein said step of signing includes the steps of

10 defining a plurality of  $k$  sub-tasks in accordance with

11  $C_1 \equiv M_1^{d_1} \pmod{p_1},$

12  $C_2 \equiv M_2^{d_2} \pmod{p_2},$

13  $\vdots$

14  $C_k \equiv M_k^{d_k} \pmod{p_k},$

15 where

16  $M_1 \equiv M \pmod{p_1},$

17  $M_2 \equiv M \pmod{p_2},$

18  $\vdots$

19  $M_k \equiv M \pmod{p_k},$

21  $d_1 \equiv d \pmod{(p_1 - 1)},$

22  $d_2 \equiv d \pmod{(p_2 - 1)}, \text{ and}$

23  $\vdots$

24  $d_k \equiv d \pmod{(p_k - 1)},$

25 wherein  $d$  is defined by

26  $d \equiv e^{-1} \pmod{(p_1 - 1) \cdot (p_2 - 1) \cdot \dots \cdot (p_k - 1)}, \text{ and}$

27  $e$  is a number relatively prime to  $(p_1 - 1), (p_2 - 1), \dots, \text{ and } (p_k - 1),$

28 solving said sub-tasks to determine results  $C_1, C_2, \dots, C_k$ , and

29 combining said results of said sub-tasks to produce said ciphertext word  $C$ .

1 52. (Once Amended) A digital signature generation system comprising:

2 a communication medium;

3 [a] digital signature generating means coupled to said communication medium and

4 operative to transform a transmit message word  $M$  to a signed ciphertext word  $C$ , and to transmit  
5 said signed ciphertext word  $C$  on said medium, wherein  $M$  corresponds to a number  
6 representative of a message, and

7  $0 \leq M \leq n-1,$

8 n being a composite number formed from the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$  wherein k is an integer  
9 greater than 2[,] and  $p_1, p_2, \dots, p_k$ , are distinct random prime numbers, and wherein the signed  
10 ciphertext word C is a number representative of a signed form of said message word M, wherein

11  $C \equiv M^d \pmod{n}$ ,

12 said digital signature generating means being operative to transform said transmit  
13 message word M to said signed ciphertext word C by performing a digital signature generating  
14 process comprising the steps of,

15 defining a plurality of k sub-tasks in accordance with,

16  $C_1 \equiv M_1^{d_1} \pmod{p_1}$ ,

17  $C_2 \equiv M_2^{d_2} \pmod{p_2}$ ,

18  $\vdots$

19  $C_k \equiv M_k^{d_k} \pmod{p_k}$ ,

20 where,

21  $M_1 \equiv M \pmod{p_1}$ ,

22  $M_2 \equiv M \pmod{p_2}$ ,

23  $\vdots$

24  $M_k \equiv M \pmod{p_k}$ ,

25

26  $d_1 \equiv d \pmod{(p_1 - 1)}$ ,

27  $d_2 \equiv d \pmod{(p_2 - 1)}$ , and

28  $\vdots$

29  $d_k \equiv d \pmod{(p_k - 1)}$ ,

30 wherein d is defined by,

31  $d \equiv e^{-1} \pmod{(p_1 - 1) \cdot (p_2 - 1) \cdot \dots \cdot (p_k - 1)}$ , and

32 e is a number relatively prime to  $(p_1 - 1), (p_2 - 1), \dots, (p_k - 1)$ ,

33 solving said sub-tasks to determine results  $C_1, C_2, \dots, C_k$ , and

34 combining said results of said sub-tasks to produce said signed ciphertext word C.

1 57. (Once Amended) A digital signature process comprising the steps of:  
2 signing a plaintext message word M to create a signed ciphertext word C, wherein M  
3 corresponds to a number representative of a message and wherein

4  $0 \leq M \leq n-1$

5 wherein n is a composite number formed by the product of  $p_1 \cdot p_2 \cdot \dots \cdot p_k$ , k is an integer  
6 greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime numbers, C is a number  
7 representative of a signed form of message word M, and wherein said encoding step  
8 comprises transforming said message word M to said ciphertext word C whereby,

9  $C \equiv M^d \pmod{n}$ ,

10 wherein d is defined by

11  $d \equiv e^{-1} \pmod{(p_1-1) \cdot (p_2-1) \cdot \dots \cdot (p_k-1)}$ , and

12 e is a number relatively prime to  $(p_1-1), (p_2-1), \dots, (p_k-1)$ ; and

13 verifying said ciphertext word C to a receive message word M' by performing the steps  
14 of,

15 defining a plurality of k sub-tasks in accordance with

16  $M_1' \equiv C_1^{e_1} \pmod{p_1}$ ,

17  $M_2' \equiv C_2^{e_2} \pmod{p_2}$ ,

18 :

19  $M_k' \equiv C_k^{e_k} \pmod{p_k}$ ,

20 wherein

21  $C_1 \equiv C \pmod{p_1}$ ,

22  $C_2 \equiv C \pmod{p_2}$ ,

23 :

24  $C_k \equiv C \pmod{p_k}$ ,

26  $e_1 \equiv e \pmod{(p_1-1)}$ ,

27  $e_2 \equiv e \pmod{(p_2-1)}$ , and

28 :

29  $e_k \equiv e \pmod{(p_k-1)}$ ,

30 solving said sub-tasks to determine results  $M_1'$ ,  $M_2'$ , ...,  $M_k'$ , and  
31 combining said results of said sub-tasks to produce said receive message word  
32  $M'$ , whereby  $M' = M$ .

1 62. (Once Amended) A digital signature system comprising:  
2 a communication medium;  
3 [a] digital signature generating means coupled to said communication medium and  
4 adapted for transforming a message word  $M$  to a signed ciphertext word  $C$  and for transmitting  
5 said signed ciphertext word  $C$  on said medium, wherein  $M$  corresponds to a number  
6 representative of a message, and  
7  $0 \leq M \leq n-1$ , wherein  $n$  is a composite number of the form  
8  $n = p_1 \cdot p_2 \cdot \dots \cdot p_k$ ,  
9 wherein  $k$  is an integer greater than 2[,] and  $p_1, p_2, \dots, p_k$  are distinct random prime  
10 numbers, and wherein said signed ciphertext word  $C$  corresponds to a number representative of a  
11 signed form of said message word  $M$  and corresponds to  
12  $C \equiv M^d \pmod{n}$ ,  
13 wherein  $d$  is defined by  
14  $d \equiv e^{-1} \pmod{(p_1 - 1) \cdot (p_2 - 1) \cdot \dots \cdot (p_k - 1)}$ , and  
15  $e$  is a number relatively prime to  $(p_1 - 1)$ ,  $(p_2 - 1)$ , ..., and  $(p_k - 1)$ ; and  
16 [a] digital signature verification means communicatively coupled with said  
17 communication medium for receiving said signed ciphertext word  $C$  via said medium, and being  
18 operative to verify said signed ciphertext word  $C$  by performing the steps of,  
19 defining a plurality of  $k$  sub-tasks in accordance with  
20  $M_1' \equiv C_1^{e_1} \pmod{p_1}$ ,  
21  $M_2' \equiv C_2^{e_2} \pmod{p_2}$ ,  
22  $\vdots$   
23  $M_k' \equiv C_k^{e_k} \pmod{p_k}$ ,  
24 wherein  
25  $C_1 \equiv C \pmod{p_1}$ ,

26  $C_2 \equiv C \pmod{p_2},$

27  $\vdots$

28  $C_k \equiv C \pmod{p_k},$

29

30  $e_1 \equiv e \pmod{(p_1 - 1)},$

31  $e_2 \equiv e \pmod{(p_2 - 1)},$

32  $\vdots$

33  $e_k \equiv e \pmod{(p_k - 1)},$

34 solving said sub-tasks to determine results  $M_1', M_2', \dots M_k'$ , and  
35 combining said results of said sub-tasks to produce said receive message word  $M'$

36 [whereby] wherein  $M' = M$ .

1 67. (New) A method as recited in claim 14 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 68. (New) A method as recited in claim 14 wherein each of said distinct random prime  
2 number has the same number of bits.

1 69. (New) A method as recited in claim 15 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 70. (New) A method as recited in claim 15 wherein each of said distinct random prime  
2 number has the same number of bits.

1 71. (New) A cryptographic communications system as recited in claim 16 wherein said step  
2 of solving said sub-tasks includes processing each of said sub-tasks by an associated one of a  
3 plurality of exponentiator units operating substantially simultaneously.

1 72. (New) A cryptographic communications system as recited in claim 16 wherein each of  
2 said distinct random prime number has the same number of bits.

1 73. (New) A method as recited in claim 17 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 74. (New) A method as recited in claim 17 wherein each of said distinct random prime  
2 number has the same number of bits.

1 75. (New) A cryptographic communications system as recited in claim 22 wherein said step  
2 of solving said sub-tasks includes processing each of said sub-tasks by an associated one of a  
3 plurality of exponentiator units operating substantially simultaneously.

~~D~~ 1 76. (New) A cryptographic communications system as recited in claim 22 wherein each of  
2 said distinct random prime number has the same number of bits.

~~H~~ 1 77. (New) A method as recited in claim 27 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 78. (New) A method as recited in claim 27 wherein each of said distinct random prime  
2 number has the same number of bits.

1 79. (New) A cryptographic communications system as recited in claim 32 wherein said step  
2 of solving said sub-tasks includes processing each of said sub-tasks by an associated one of a  
3 plurality of exponentiator units operating substantially simultaneously.

1 80. (New) A cryptographic communications system as recited in claim 32 wherein each of  
2 said distinct random prime number has the same number of bits.

1 81. (New) A method as recited in claim 37 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 82. (New) A method as recited in claim 37 wherein each of said distinct random prime  
2 number has the same number of bits.

1 83. (New) A cryptographic communications system as recited in claim 42 wherein said step  
2 of solving said sub-tasks includes processing each of said sub-tasks by an associated one of a  
3 plurality of exponentiator units operating substantially simultaneously.

1 84. (New) A cryptographic communications system as recited in claim 42 wherein each of  
2 said distinct random prime number has the same number of bits.

1 85. (New) A method as recited in claim 47 wherein said step of solving said sub-tasks  
2 includes processing each of said sub-tasks by an associated one of a plurality of exponentiator  
3 units operating substantially simultaneously.

1 86. (New) A method as recited in claim 47 wherein each of said distinct random prime  
2 number has the same number of bits.

1 87. (New) A digital signature generation system as recited in claim 52 wherein said step of  
2 solving said sub-tasks includes processing each of said sub-tasks by an associated one of a  
3 plurality of exponentiator units operating substantially simultaneously.

1 88. (New) A digital signature generation system as recited in claim 52 wherein each of said  
2 distinct random prime number has the same number of bits.

1 89. (New) A digital signature process as recited in claim 57 wherein said step of solving said  
2 sub-tasks includes processing each of said sub-tasks by an associated one of a plurality of  
3 exponentiator units operating substantially simultaneously.